

# kNN Decision Boundary at $k=1$ and Voronoi Tessellation

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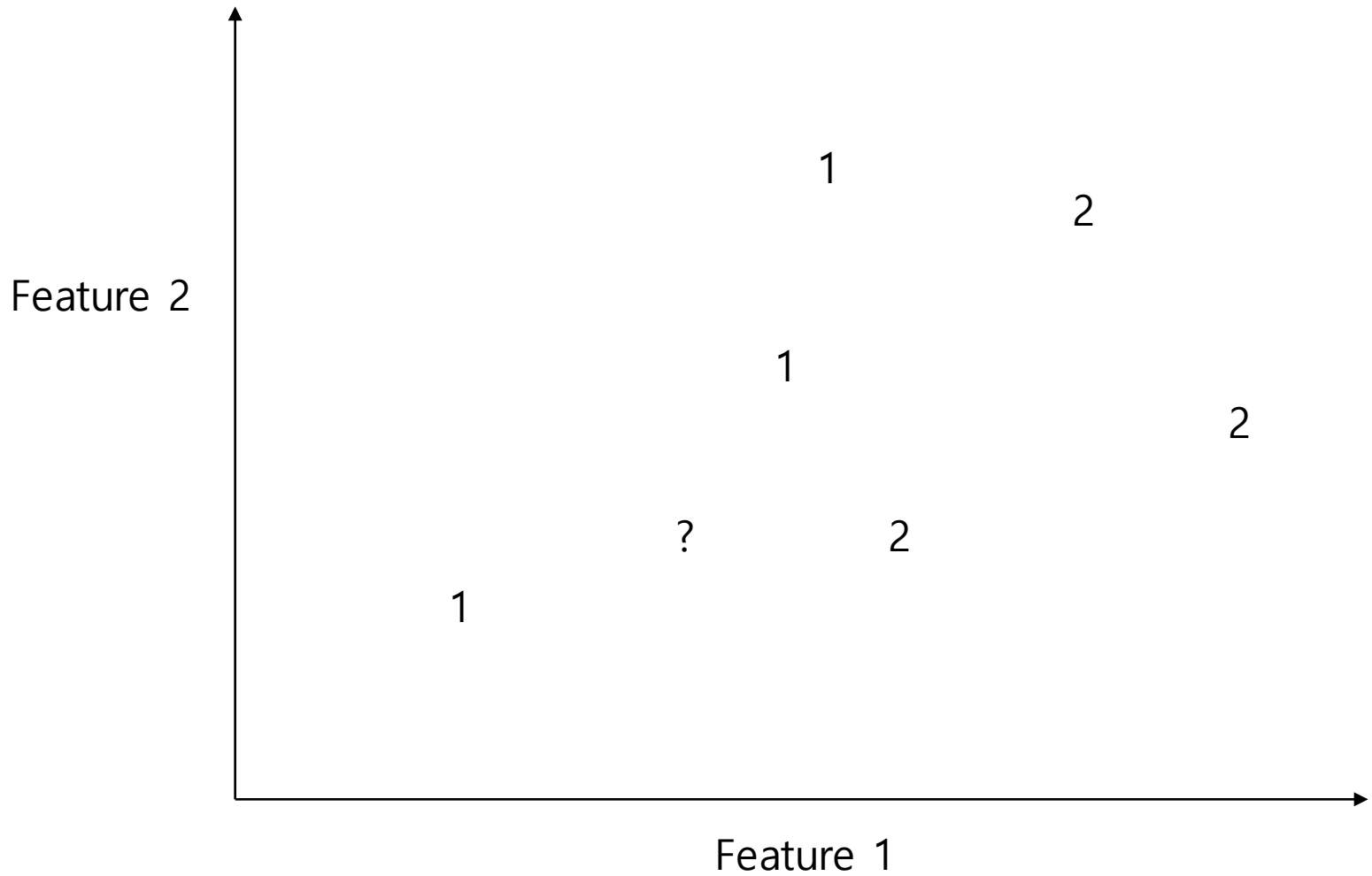
<http://link.koreatech.ac.kr>

# kNN Algorithm at $k=1$

## ◆ Algorithm

- $y$  is a new feature vector whose class label is unknown
- Search the training set for the closest feature vector to  $y$ 
  - let this “closest feature vector” be  $x(j)$
  - How are “closest  $x$ ” vectors determined?
    - typically use minimum Euclidean distance
- Classify  $y$  with the same label as  $x(j)$ ,
  - i.e.  $y$  is assigned label  $c(j)$
- Side note: this produces a “Voronoi tessellation” of the  $d$ -space
  - each point “claims” a cell surrounding it
  - cell boundaries are polygons

# How should the new point be classified?

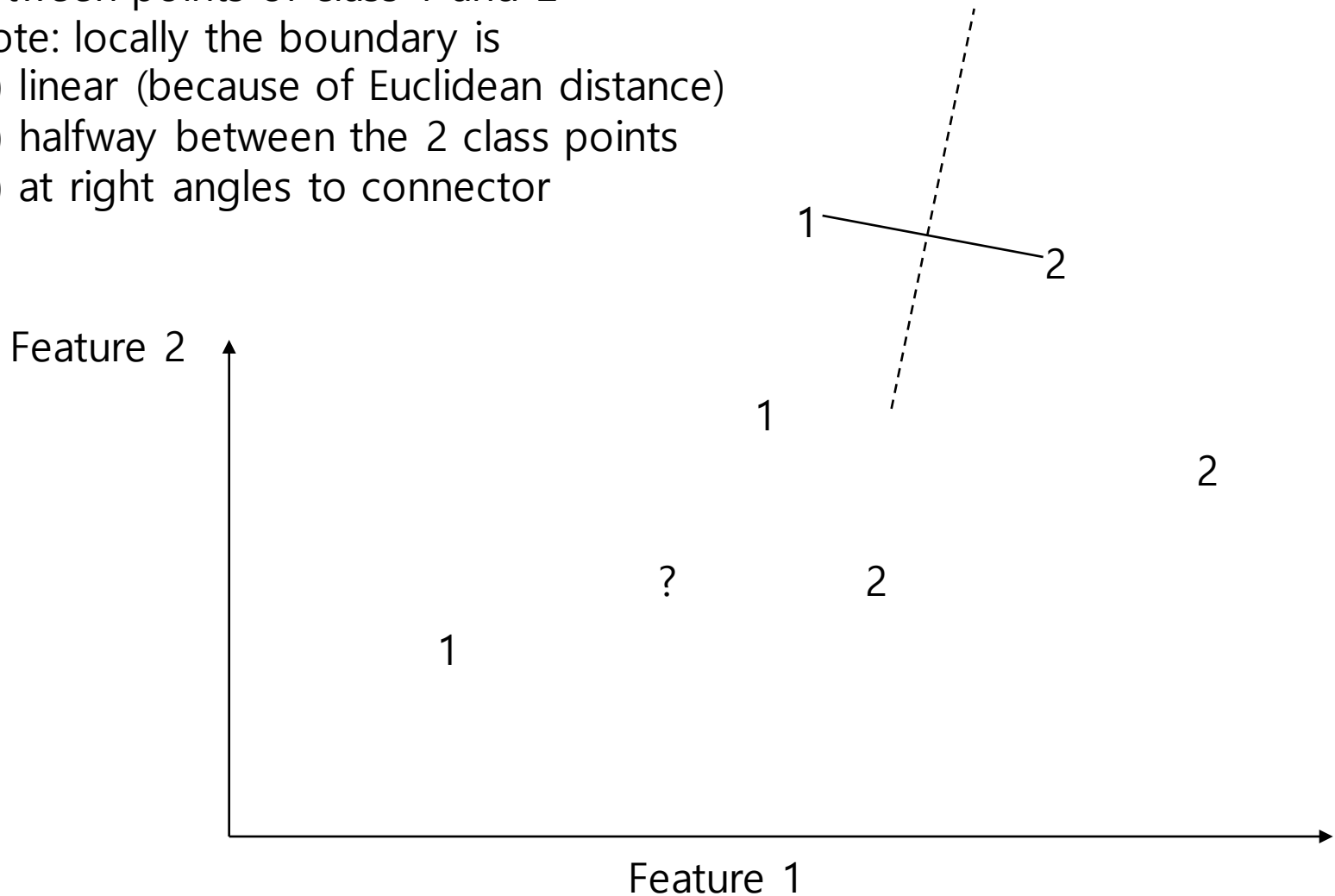


# Local Decision Boundaries

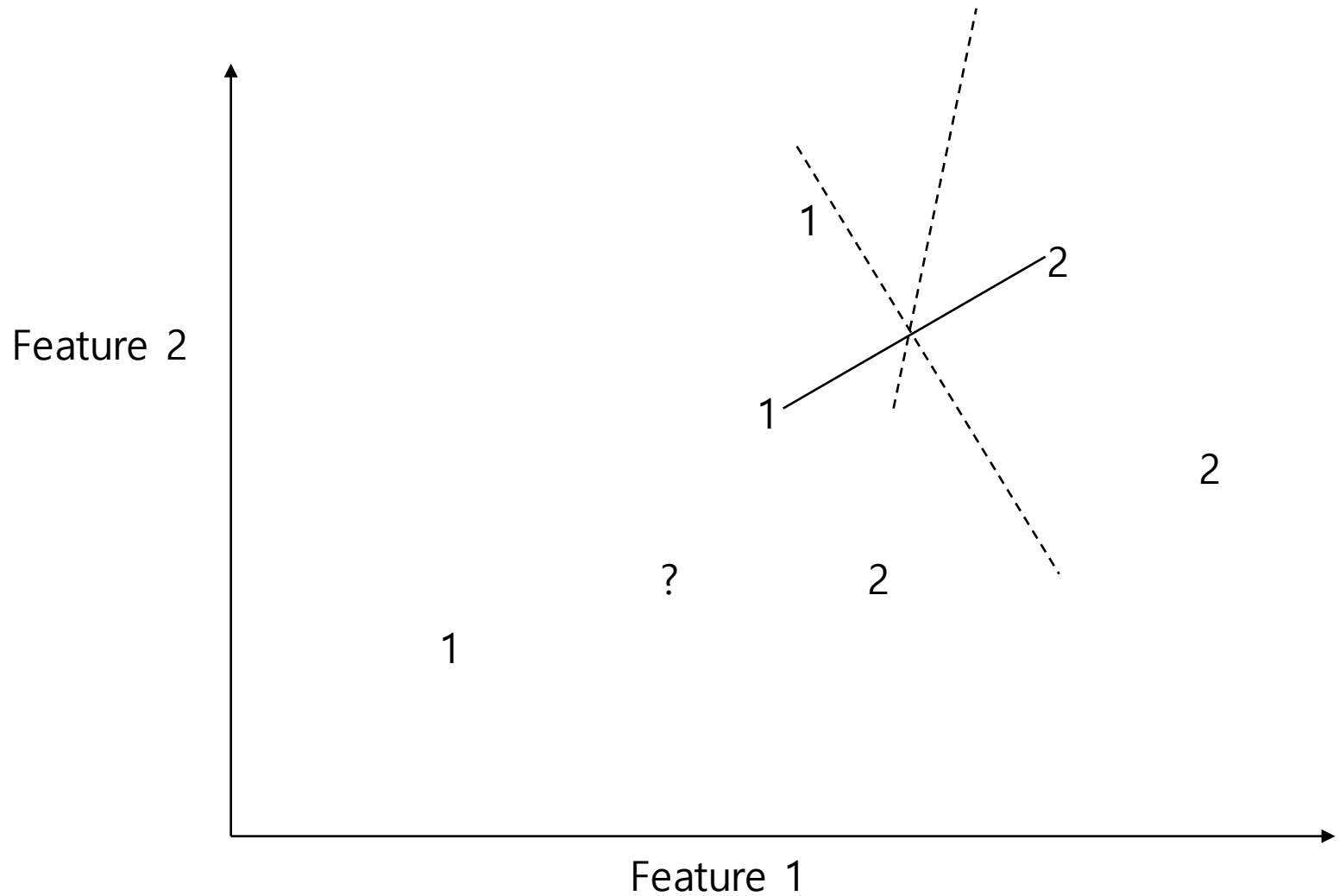
Boundary? Points that are equidistant between points of class 1 and 2

Note: locally the boundary is

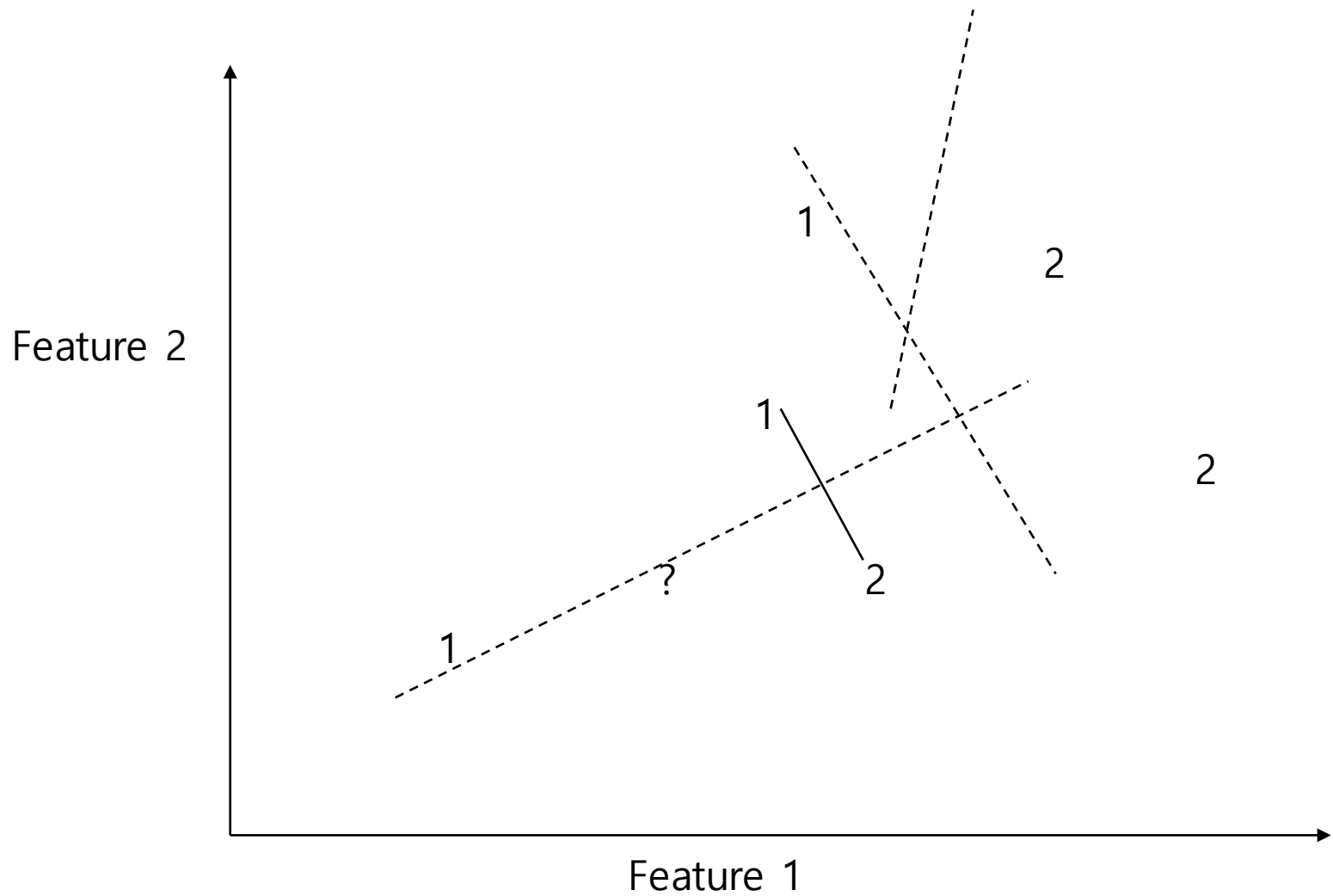
- (1) linear (because of Euclidean distance)
- (2) halfway between the 2 class points
- (3) at right angles to connector



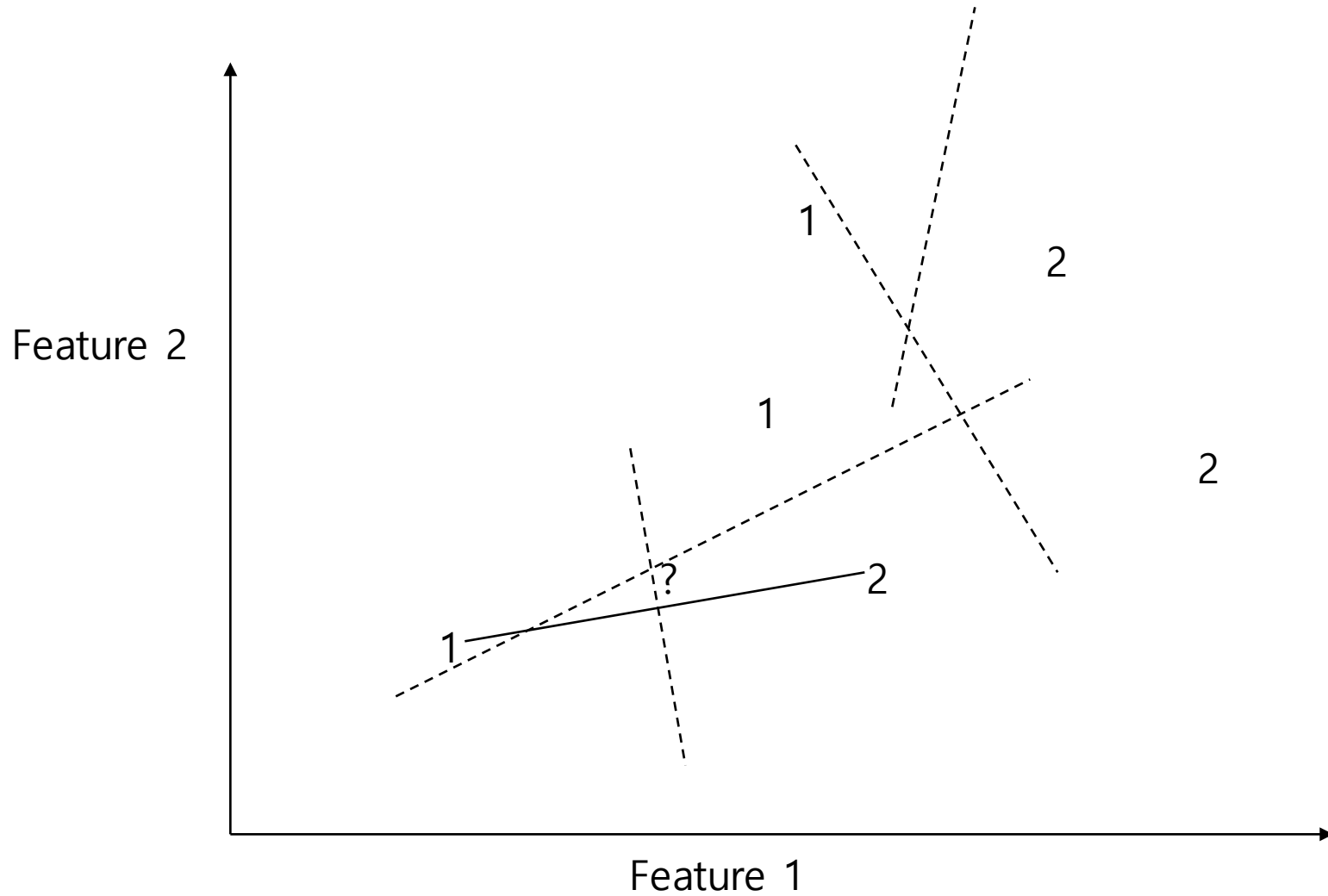
# Finding the Decision Boundaries



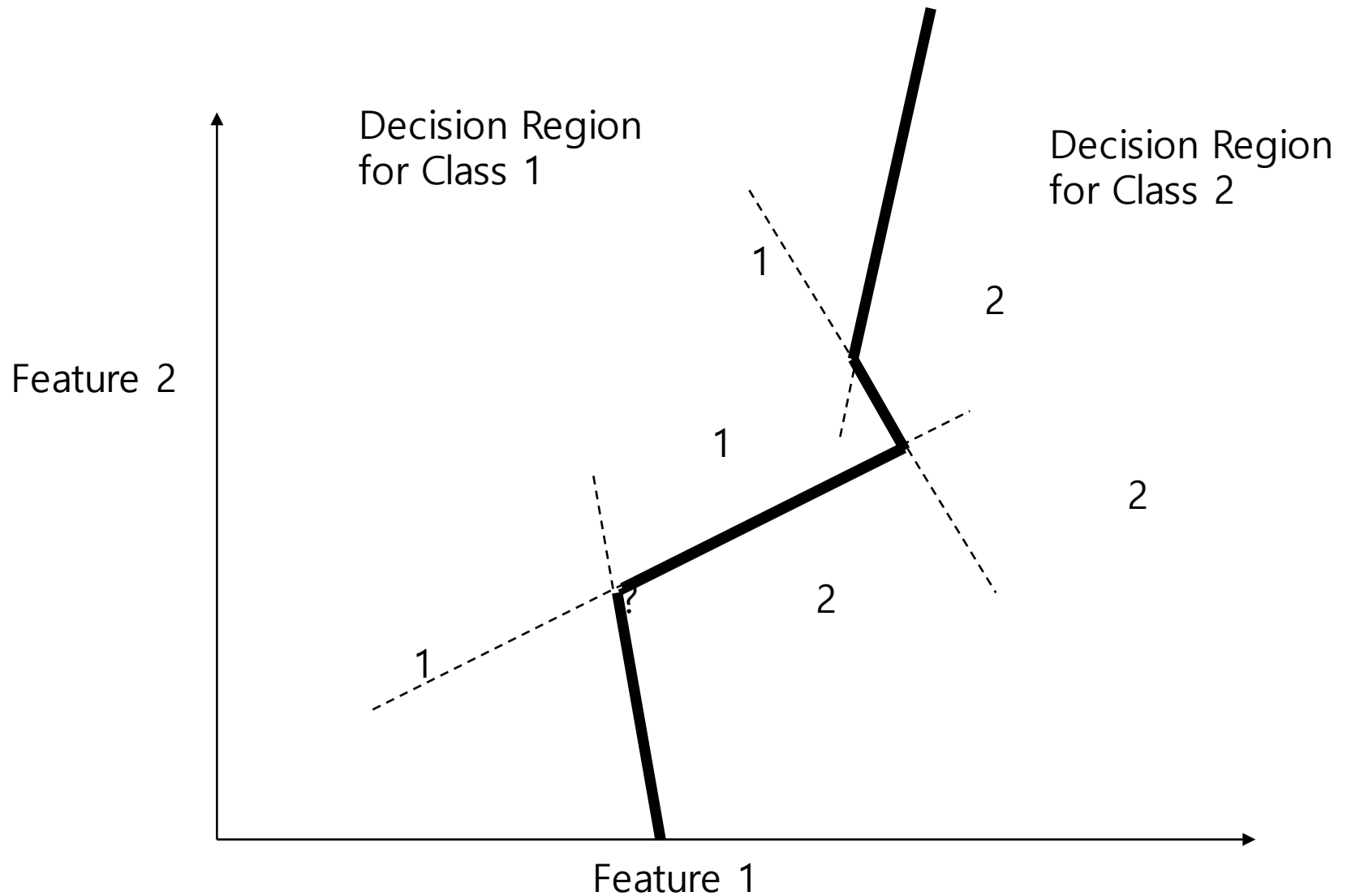
# Finding the Decision Boundaries



# Finding the Decision Boundaries

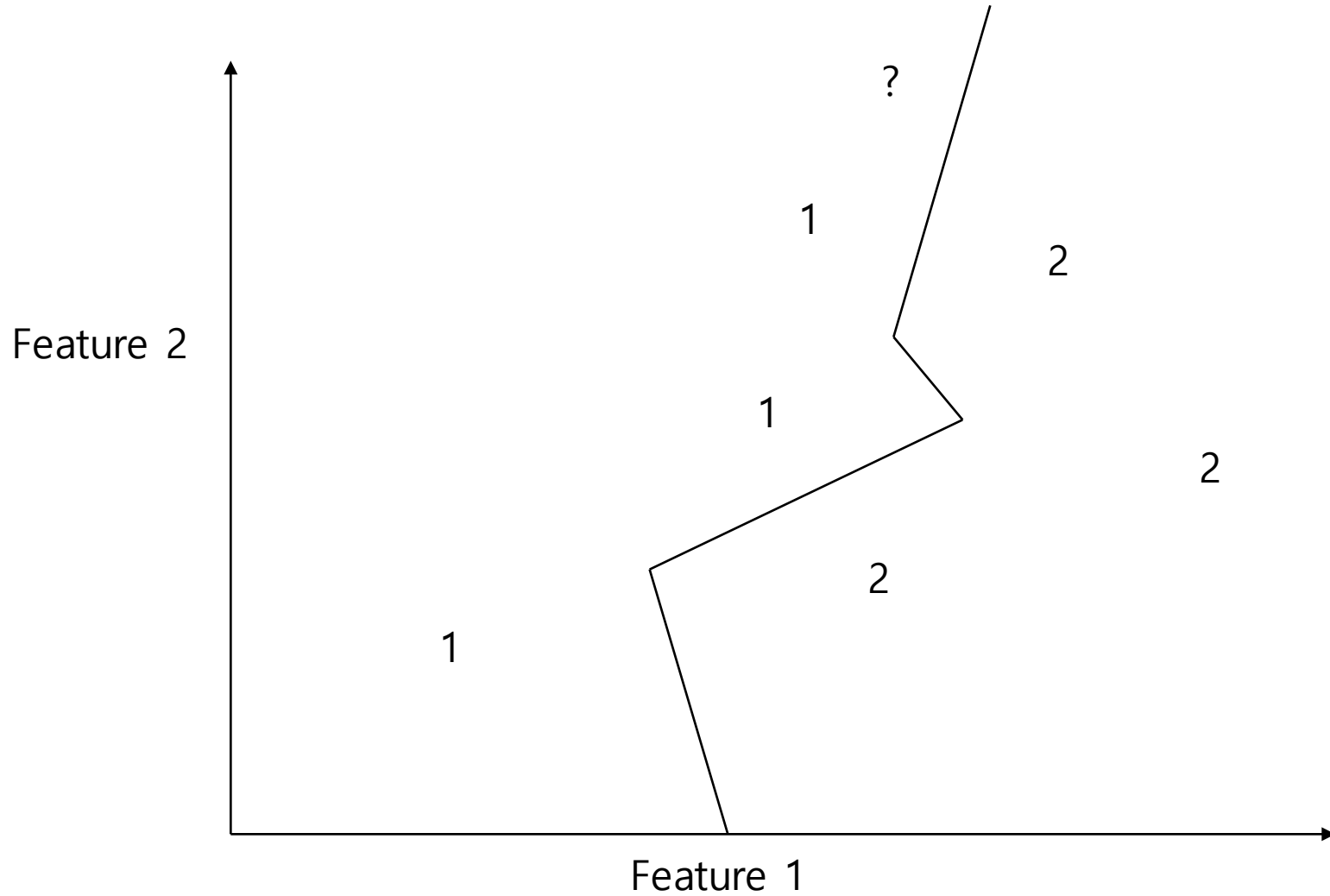


# Overall Boundary = Piecewise Linear

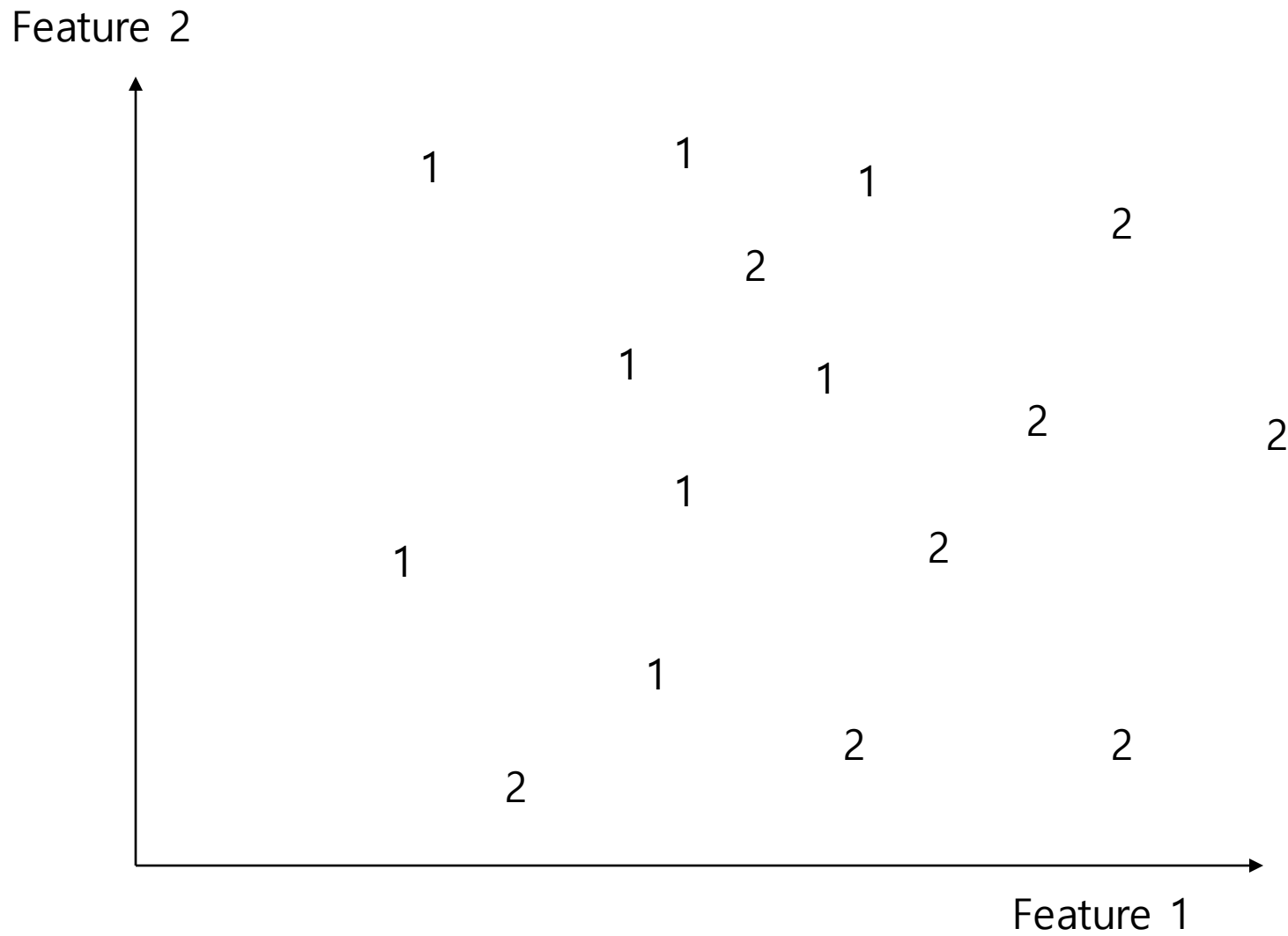




# Geometric Interpretation of kNN (k=1)



# More Data Points



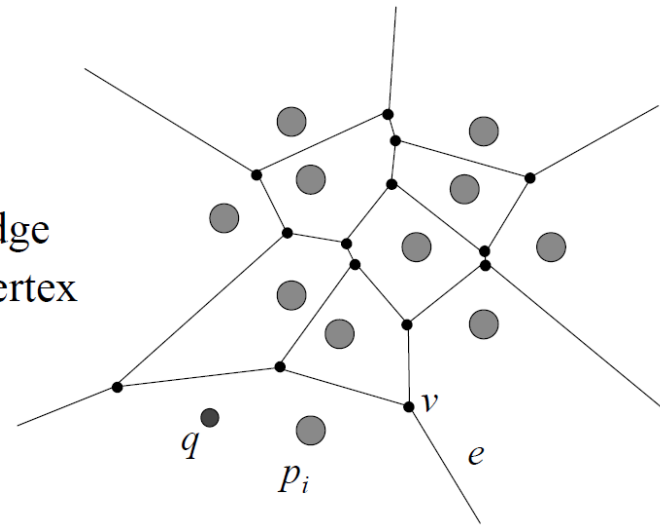


# Voronoi diagram (tesellation)

## ◆ Voronoi diagram

- The Voronoi diagram is formed from lines that bisect and are perpendicular to the lines that connect two neighboring vertices.
- Each point  $s$  has a **Voronoi cell**  $V(s)$  consisting of all points closer to  $s$  than to any other points

$p_i$  : site points  
 $q$  : free point  
 $e$  : Voronoi edge  
 $v$  : Voronoi vertex



A point  $q$  lies in the cell corresponding to a site  $p_i \in P$   
*iff*

$\text{Euclidean\_Distance}(q, p_i) < \text{Euclidean\_distance}(q, p_j)$ ,  
for each  $p_i \in P, j \neq i$ .

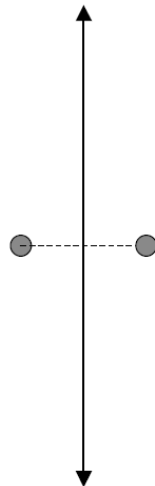
# Voronoi diagram

## ◇ Voronoi diagram examples

– 1 point



– 2 points form “a perpendicular bisector”

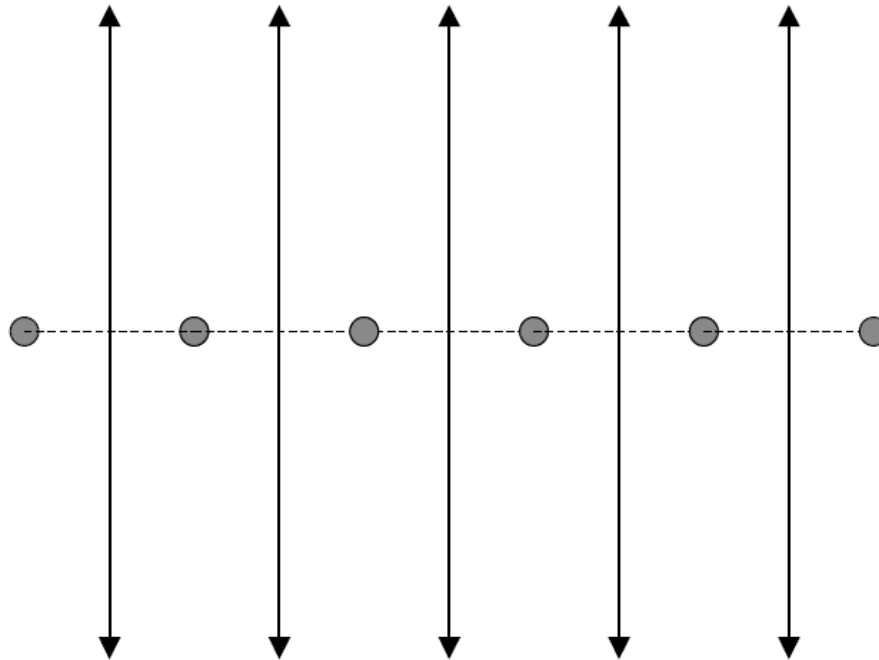


Voronoi Diagram is a line that extends infinitely in both directions, and the two half planes on either side.

# Voronoi diagram

## ◇ Voronoi diagram examples

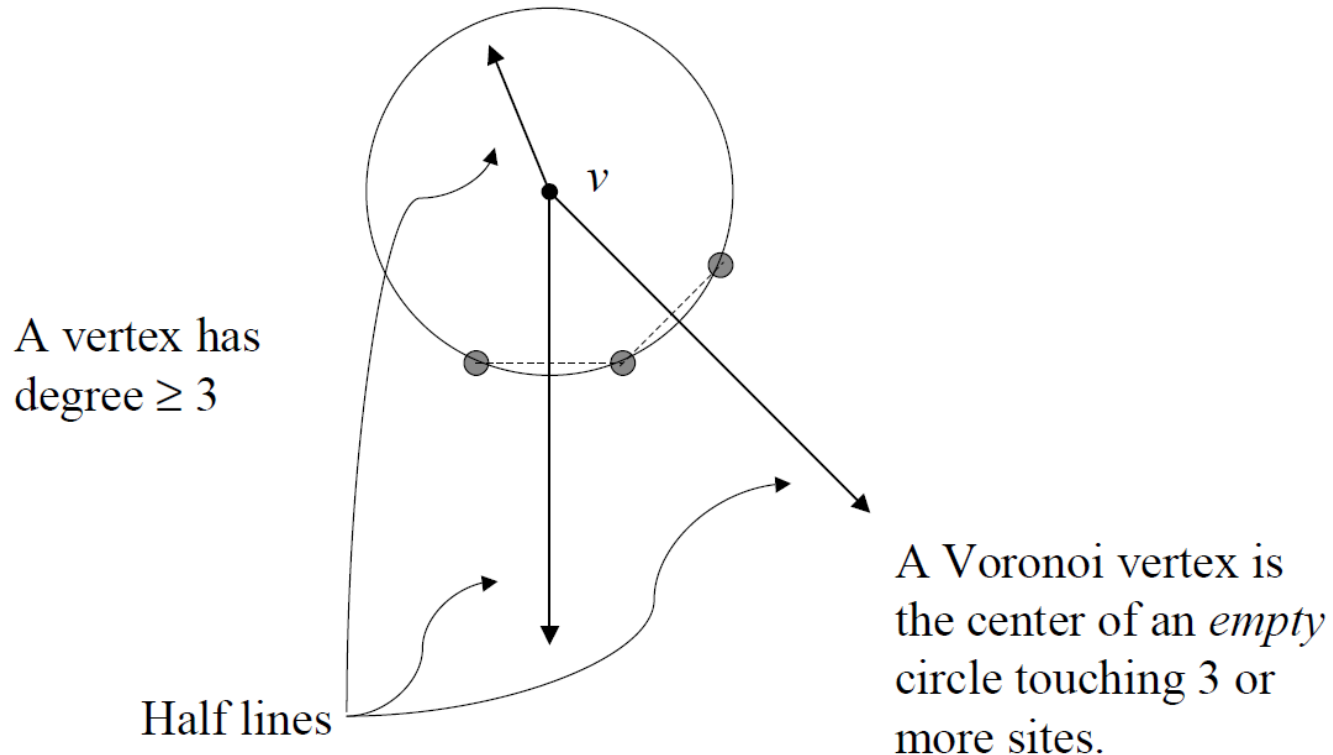
- Collinear points form “a series of parallel lines”



# Voronoi diagram

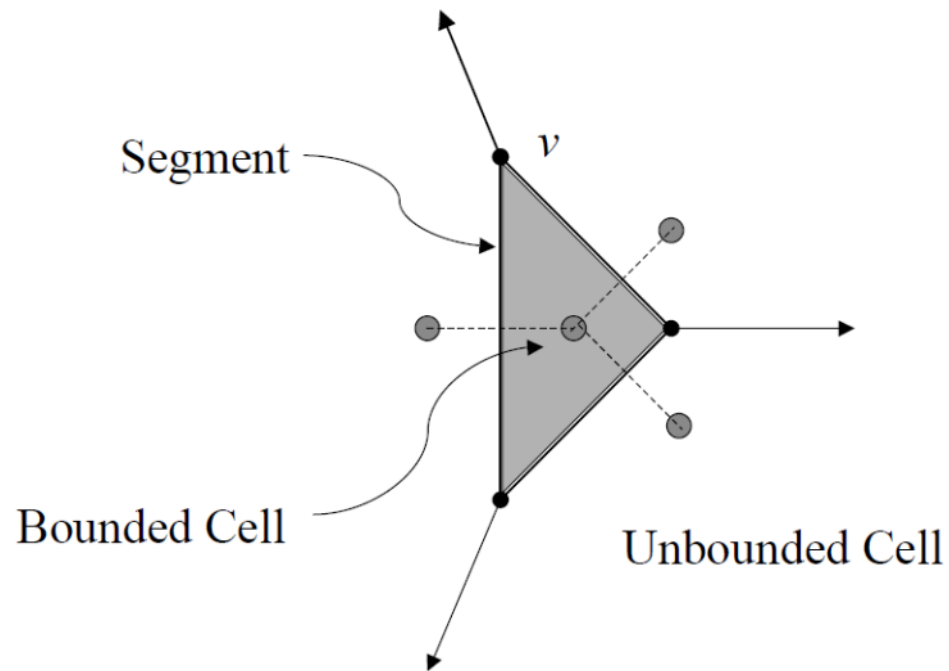
## ◆ Voronoi diagram examples

- Non-collinear points form “Voronoi half lines” that meet at a vertex



# Voronoi diagram

## ◇ Voronoi cells and segments



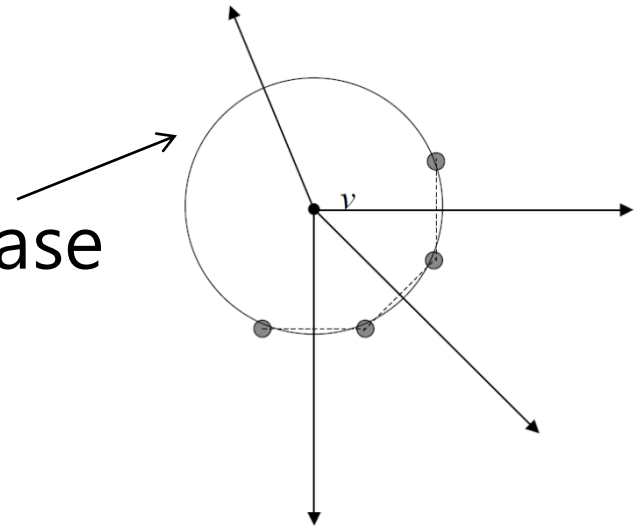


# Voronoi diagram

◆ Which of the following is true for 2-D Voronoi diagrams?

- Four or more non-collinear points are...
  - 1) sufficient to create a bounded cell
  - 2) necessary to create a bounded cell
  - 3) 1 and 2
  - 4) none of above

◆ Four points' degeneration case of bounded cell:

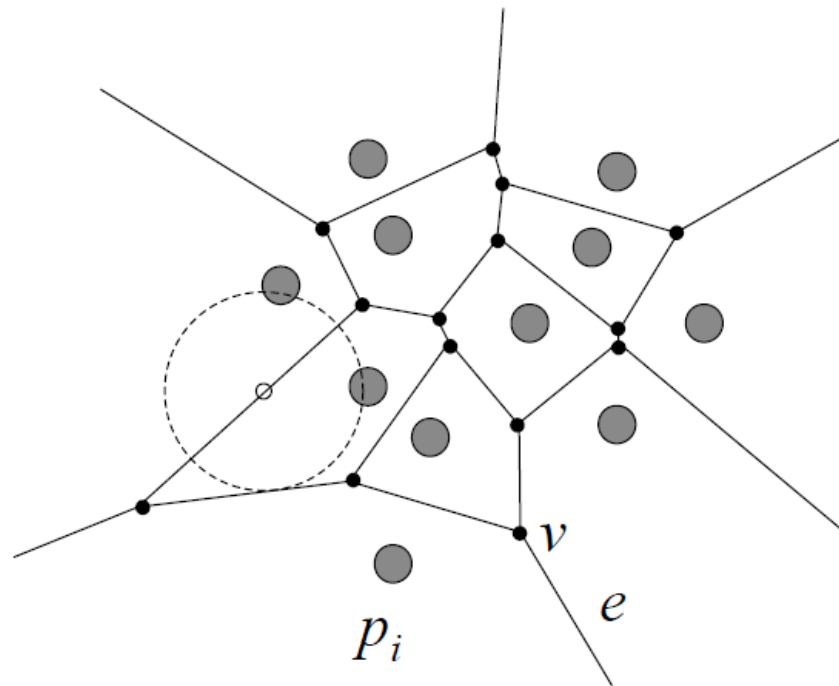


# Property I of Voronoi diagram

A point  $q$  lies on a Voronoi edge between sites  $p_i$  and  $p_j$  iff the largest empty circle centered at  $q$  touches only  $p_i$  and  $p_j$

- A Voronoi edge is a subset of locus of points equidistant from  $p_i$  and  $p_j$

$p_i$  : site points  
 $e$  : Voronoi edge  
 $v$  : Voronoi vertex

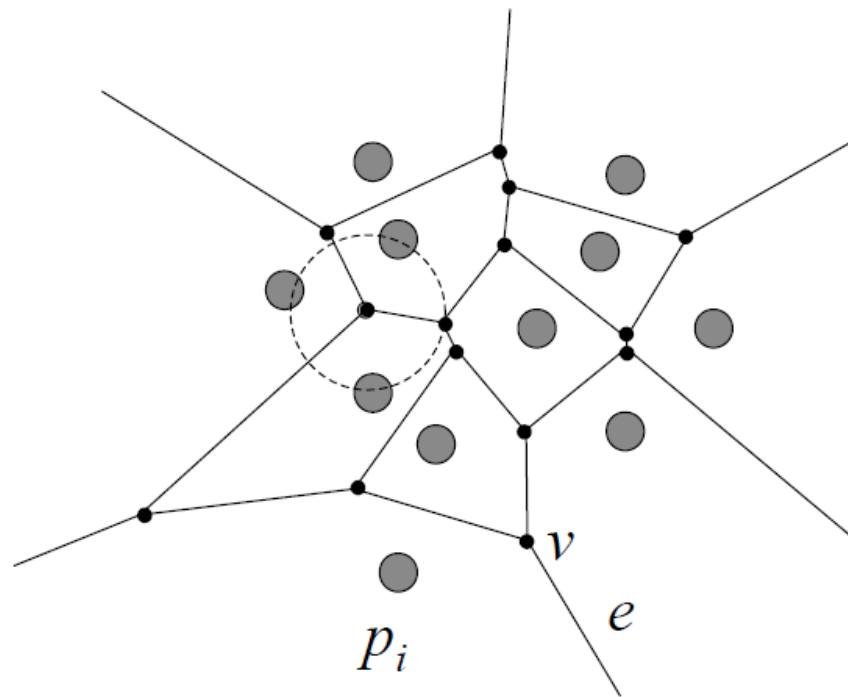


# Property II of Voronoi diagram

A point  $q$  is a vertex *iff* the largest empty circle centered at  $q$  touches at least 3 sites

- A Voronoi vertex is an intersection of 3 more segments, each equidistant from a pair of sites

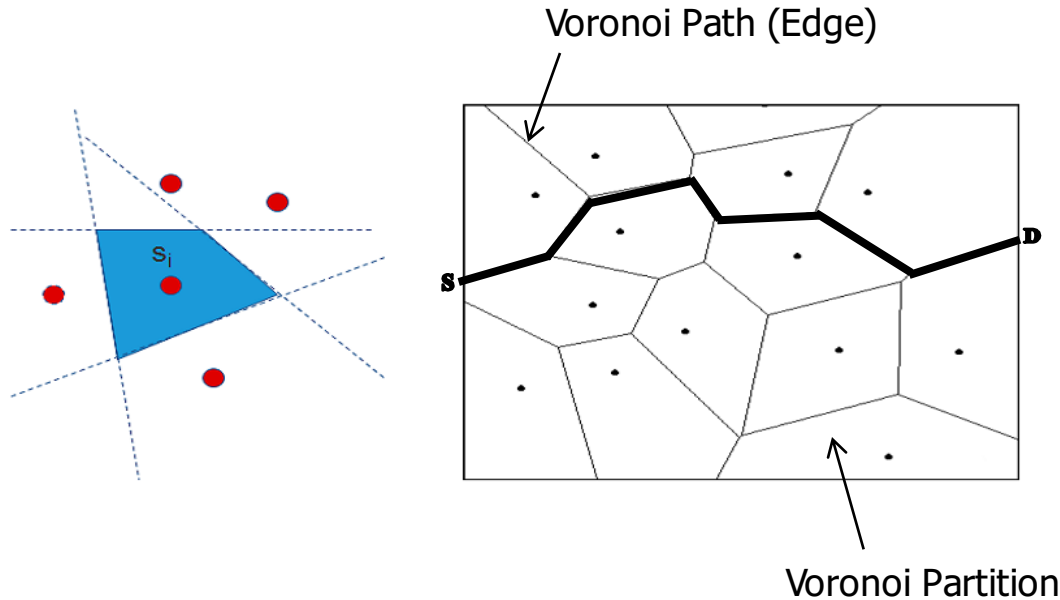
$p_i$  : site points  
 $e$  : Voronoi edge  
 $v$  : Voronoi vertex



# Surveillance Path

## ◇ Maximal Breach Path

Voronoi diagram



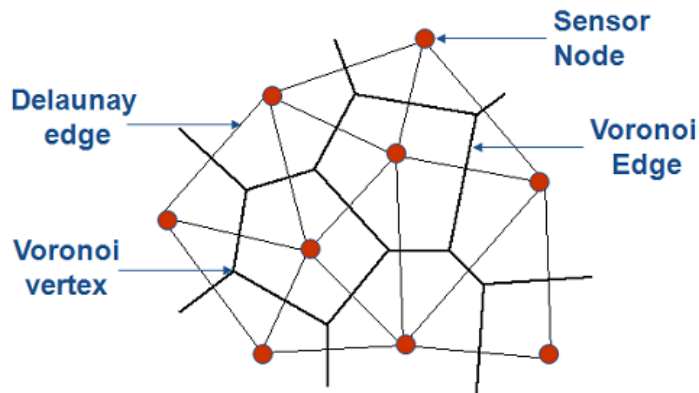
### – Voronoi Path (= Maximal Breach Path)

- The path where the surveillance level is the lowest
- The path where its closest distance to any sensor is as large as possible.

# Surveillance Path

## ◇ Maximal Support Path

Delaunay triangulation



## – Delaunay Triangulation Path (= Maximal Support Path)

- The path where the surveillance level is the greatest
- The path where its closest distance to any sensor is as short as possible.